



Coding of
substitution
dynamical
systems

Paul Surer

Motivation

Desubstitution

Coding

Applications

Coding of substitution dynamical systems

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unesp 

 **FAPESP**

Relation between beta-expansions and substitutions



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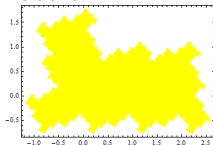
Motivation

Desubstitution

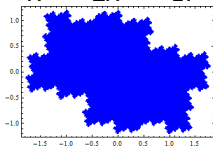
Coding

Applications

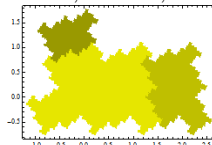
Central tile induced by the
beta-expansion with β the
real root of $x^3 - 2x^2 - 1$.



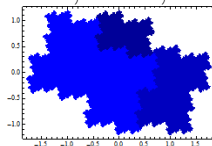
Central tile induced by the
symmetric beta-expansion
with β the real root of
 $x^3 - 2x^2 - 1$.



Rauzy fractal associated
with the substitution
 $1 \mapsto 112, 2 \mapsto 3, 3 \mapsto 1$.



Rauzy fractal associated
with the substitution
 $1 \mapsto 121, 2 \mapsto 3, 3 \mapsto 1$.





Possible Explanations

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- ~~Paul messed up with Mathematica[®].~~
- ~~There is a accidental similarity.~~
- There is an ~~unknown~~ relation between the shift of finite type induced by the symmetric beta expansion and the substitution dynamical system.



A basic result

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Theorem (cf. Mossé, Durand et al., Canterini-Siegel, Holton-Zamboni, ...)

For every $w \in \Omega$ there is a unique pair $(u, k) \in \Omega \times \mathbb{N}$ with $0 \leq k < |\sigma(u_0)|$ such that $w = S^k \sigma(u)$.

Corollary

For every $w \in \Omega$ there is a unique pair $(u', k') \in \Omega \times \mathbb{Z}$ with $-|\sigma(u_{-1})| < k' \leq 0$ such that $w = S^{k'} \sigma(u)$.

Idea

We combine these two assertions. To maintain uniqueness we use a coding prescription.



Coding prescription

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Definition (Coding prescription)

A coding prescription is a collection $(\mathcal{C}(j))_{j \in \mathcal{A}}$ where each $\mathcal{C}(j)$ is a complete set of residues modulo $|\sigma(j)|$ whose entries are smaller than $|\sigma(j)|$ in modulus.

Notation

For a coding prescription \mathcal{C} we denote for each $j \in \mathcal{A}$ the non-negative elements of $\mathcal{C}(j)$ by $\mathcal{C}(j)^+$ and the non-positive elements by $\mathcal{C}(j)^-$.

Theorem

Let \mathcal{C} a coding prescription. Then for every $w \in \Omega$ there is a unique pair $(\Theta(w), \phi(w)) \in \Omega \times \mathbb{Z}$ such that $\phi(w) \in \mathcal{C}(\Theta(w)_{-1})^- \cup \mathcal{C}(\Theta(w)_0)^+$ and $w = S^{\phi(w)}\sigma(\Theta(w))$.



Inverse letters

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Notation

Let $\tilde{\mathcal{A}} := \mathcal{A} \cup \{\bar{j} \mid j \in \mathcal{A}\}$ and $\tilde{\mathcal{A}}^*$ the finite words over $\tilde{\mathcal{A}}$.

Rules

Observe the following rules and notations: for all $A, B \in \tilde{\mathcal{A}}^*$ we have

- $\overline{AB} = \overline{B} \overline{A}$;
- $\overline{(\overline{A})} = A$ for all $A \in \tilde{\mathcal{A}}^*$;
- $A\overline{A} = \varepsilon$ for all $A \in \tilde{\mathcal{A}}^*$;
- $\overline{\varepsilon} = \varepsilon$;
- $\sigma(\overline{A}) = \overline{\sigma(A)}$.

"Digits expansion" for substitutions

Digit set

$$\mathcal{R} = \mathcal{R}(\mathcal{C}) := \{(a.b, d) \mid ab \in \mathcal{L}_2, \exists k \in \mathcal{C}(a)^- \cup \mathcal{C}(b)^+, \sigma(ab)_{[1, |\sigma(a)|+k]} = \sigma(a)d\}.$$

Expansion

$\Gamma : \Omega \longrightarrow \mathcal{R}^\infty, w \mapsto (\Theta^n(w)_{-1} \cdot \Theta^n(w)_0, d_n)_{n \geq 1}$
where, for every $n \geq 1$, d_n is the (uniquely determined) element of $\tilde{\mathcal{A}}^*$ with $\Theta^{n-1}(w) = \sigma(\Theta^n(w)_{(-\infty, -1]})d_n \cdot \overline{d_n} \sigma(\Theta^n(w)_{[0, \infty)})$.

Note

- d_n consists of negative or positive letters only for all $n \in \mathbb{N}$.
- With the convention that inverse letters give negative contributions to the length of a word we have $|d_n| = \phi(\Theta^{n-1}(w))$.





Properties

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Proposition

Let $\Gamma(w) = (a_n \cdot b_n, d_n)_{n \geq 1}$. If for infinitely many indices n we have $|\sigma(a_n)d_n| > 1$ and $|\overline{d_n}\sigma(b_n)| > 1$, respectively then $w = \lim_{n \rightarrow \infty} \sigma^n(a_n)\sigma^{n-1}(d_n) \cdots d_1 \cdot \overline{d_1} \cdots \sigma^{n-1}(\overline{d_n})\sigma^n(b_n)$.

Proposition

w is a periodic point if and only if $\Gamma(w) = (a_n \cdot b_n, d_n)_{n \geq 1}$ with $d_n = \varepsilon$ for all $n \geq 1$.



Graph representation

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Definition

The *coding graph with respect to \mathcal{C}* , denoted by $G(\sigma, \mathcal{C})$, is the graph with the following properties: The vertices are pairs $(a.b)$ such that $ab \in \mathcal{L}_2$. There is an edge from $(a.b)$ to $(a'.b')$ labelled by $(a'.b', d) \in \mathcal{R}$ if $(\sigma(a')d)_{|\sigma(a')d|} = a$ and $(\bar{d}\sigma(b'))_1 = b$.

Theorem

Γ maps Ω surjectively onto the infinite paths of $G(\sigma, \mathcal{C})$. Γ is one-to-one with the shift orbit of periodic points as possible exceptions (finite-to-one).



Partial ordering on the edges

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Definition

A coding prescription \mathcal{C} is called *continuous* if $\mathcal{C}(j)$ consists of consecutive integers for all $j \in \mathcal{A}$.

Partial ordering

The partial ordering on \mathcal{R}
 $(a.b, d) \prec (a'.b', d')$ if and only if $a = a', b = b', |d| < |d'|$
induces a partial ordering on the paths.

Note

For a continuous coding prescription the successor map is a "nice" adic transformation. For non-continuous coding prescription the successor map is difficult to describe.



The shift map

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Theorem

The successor map on the paths is conjugated to the shift map on Ω .

Theorem

If $\Gamma(w)$ is a maximal (minimal, respectively) element (w.r.t. \prec) then $\Gamma(S(w))$ is a minimal element ($\Gamma(S^{-1}(w))$ is a maximal element, respectively).



Symmetric beta-expansions

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Conjecture

Let $\beta > 1$ and $d(-\frac{1}{2}) = -t_1, \dots, -t_q, -t_{q+1}, \dots, -t_{q+p}$ such that $t_n \geq 0$ for all $n \in \{1, \dots, p+q\}$. Then the symmetric beta-shift is a factor of the substitution dynamical system induced by

$$1 \mapsto 1^{t_1} 2 1^{t_1},$$

$$\vdots$$

$$(p+q-1) \mapsto 1^{t_{p+q-1}} 2 1^{t_{p+q-1}},$$

$$(p+q) \mapsto 1^{t_{p+q}} (q+1) 1^{t_{p+q}}.$$

Epsilon-beta-expansions



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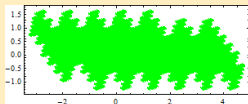
Applications

S. (2009), Kalle-Steiner (2011)

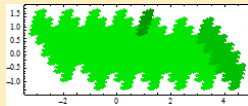
Beta-expansions of the interval $[-\varepsilon, 1 - \varepsilon)$ with respect to the transformation $x \mapsto x\beta - \lfloor x\beta + \varepsilon \rfloor$ ($\varepsilon \in [0, 1)$).

Example: $\beta^6 - 6\beta^2 - 3\beta - 1 = 0$, $\varepsilon = \frac{1}{3}$

The induces central tile



The Rauzy fractal induced
by $1 \mapsto 1111611, 2 \mapsto$
 $1131, 3 \mapsto 1$.





Epsilon-beta-expansions

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Conjecture

Let $\beta > 1$, $d(-\varepsilon) = -t_1, \dots, -t_q, -t_{q+1}, \dots, -t_{q+p}$ and $d(1 - \varepsilon) = s_1, \dots, s_q, s_{q+1}, \dots, s_{q+p}$ such that $t_n, s_n \geq 0$ for all $n \in \{1, \dots, p + q\}$. Then the symmetric beta-shift is a factor of the substitution dynamical system induced by

$$1 \mapsto 1^{t_1} 21^{s_1},$$

$$\vdots$$

$$(p + q - 1) \mapsto 1^{t_{p+q-1}} 21^{s_{p+q-1}},$$

$$(p + q) \mapsto 1^{t_{p+q}} (q + 1) 1^{s_{p+q}}.$$

Rauzy fractals



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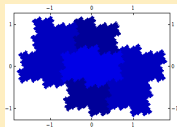
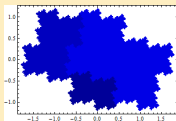
Applications

Observation

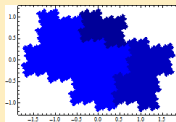
The infinite walks with initial point $a.b$ of the coding graph correspond to the intersection of the b th subtile with the a th subtile under the domain exchange map.

$$\sigma : 1 \mapsto 121, 2 \mapsto 3, 3 \mapsto 1$$

Rauzy fractal under the domain exchange



Rauzy fractal





Further ideas

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Applications

- Consider applications on Automorphisms of the free group (links to other shapes of epsilon-beta-expansions or negative beta-expansions);
- Comparing (intersecting) Rauzy fractals for different substitution that have the same incidence matrix;
- Studying tilings.



Induced substitution

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Observation

The partial ordering on the edges of the coding graph induces another primitive substitution over \mathcal{L}_2 . Thus, each continuous coding prescription \mathcal{C} gives an *induced substitution* $\sigma_{\mathcal{C}}$.

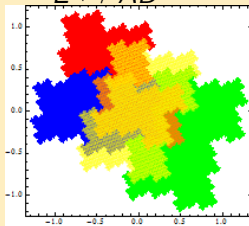
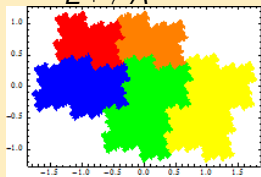
- $\sigma_{\mathcal{C}}$ has the same dominant eigenvalue as the original substitution σ ;
- $\sigma_{\mathcal{C}}$ is always reducible;
- the right eigenvector of the incidence matrix does not depend on \mathcal{C} .

Induced substitution

$$\sigma : 1 \mapsto 121, 2 \mapsto 3, 3 \mapsto 1$$

$$C_I(1) = \{-2, -1, 0\}, \quad \sigma_{C_I} : A \mapsto BCA,$$
$$B \mapsto BCD,$$
$$C \mapsto E,$$
$$D \mapsto BCA,$$
$$E \mapsto A$$
$$C_I(1) = \{-1, 0, 1\}, \quad \sigma_{C_I} : A \mapsto CAB,$$
$$B \mapsto CD,$$
$$C \mapsto EB,$$
$$D \mapsto CA,$$
$$E \mapsto AB$$

$$C_I(1) = \{0, 1, 2\}, \quad \sigma_{C_I} : A \mapsto ABC,$$
$$B \mapsto D,$$
$$C \mapsto EBC,$$
$$D \mapsto A,$$
$$E \mapsto ABC$$



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Thanks

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Thank you
for your attention

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