

SUBSTITUTIONS, SHIFTS OF FINITE TYPE, AND NUMERATION

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FWF

Der Wissenschaftsfonds.



Let $\mathcal{A} := \{1, 2, \dots, m\}$ be a finite set (alphabet) and $\mathcal{A}^* = \bigcup_{n \geq 0} \mathcal{A}^n$ the free monoid over \mathcal{A} (the empty word ε is the neutral element). For an element of \mathcal{A}^* we denote by $|A|$ the length of A .

Preliminaries

Let $\mathcal{A} := \{1, 2, \dots, m\}$ be a finite set (alphabet) and $\mathcal{A}^* = \bigcup_{n \geq 0} \mathcal{A}^n$ the free monoid over \mathcal{A} (the empty word ε is the neutral element). For an element of \mathcal{A}^* we denote by $|A|$ the length of A .

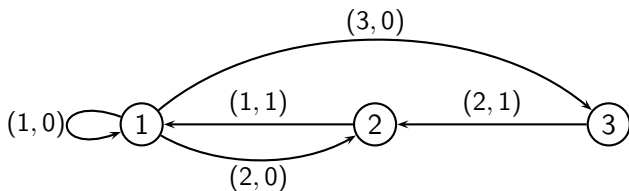
Consider a non-erasing morphism (substitution) $\sigma : \mathcal{A}^* \mapsto \mathcal{A}^*$. We call σ primitive if there exists a positive integer n such that b appears in $\sigma^n(a)$ for each pair $(a, b) \in \mathcal{A}^2$.

Graphs associated with substitutions

Substitutions are frequently represented by a finite directed graph. This graph appears under different names and in slightly different forms (labels, edge directions). For example, let $\mathcal{A} = \{1, 2, 3\}$ and $\sigma : 1 \mapsto 12, 2 \mapsto 13, 3 \mapsto 1$.

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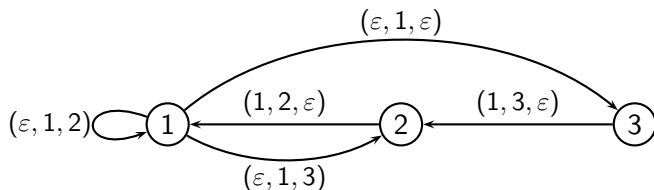
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Holton-Zambioni (2001)

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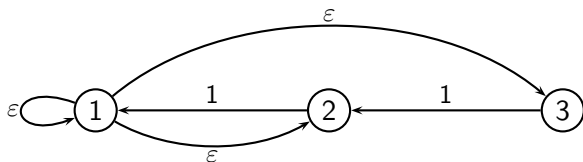
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Prefix-Suffix graph, Canterini-Siegel (2001)

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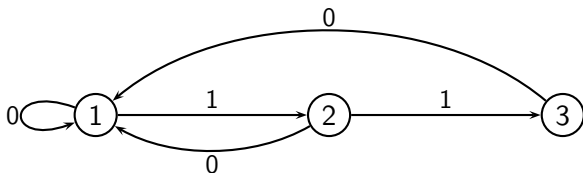
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Prefix graph (Rauzy fractals)

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Dumont-Thomas numeration

Periodic points

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with $w'_0 w'_1 w'_3 \cdots = \sigma(w_0) \sigma(w_1) \cdots$, and

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We mark this centre by a point and write

$$\sigma(\cdots w_{-2} w_{-1} \cdot w_0 w_1 w_2 w_3 \cdots) = \cdots \sigma(w_{-2}) \sigma(w_{-1}) \cdot \sigma(w_0) \sigma(w_1) \dots$$

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There exists at least one *periodic point* $u \in \mathcal{A}^{\mathbb{N}}$, that is $\sigma^n(u) = u$ for some integer $n > 1$. For avoiding trivialities we suppose that u is not periodic. Let

$$\mathfrak{L} := \{A \in \mathcal{A}^* : A \text{ is a subword of } u\}$$

the language induced by σ . Due to primitivity \mathfrak{L} does not depend on the actual choice of u .

A dynamical system

Let

$$\Omega := \{w \in \mathcal{A}^{\mathbb{Z}} : \text{each subword of } w \text{ is contained in } \mathfrak{L}\}$$

and $S : \mathcal{A}^{\mathbb{Z}} \rightarrow \mathcal{A}^{\mathbb{Z}}$ the left shift, i.e. $S : (w_j)_{j \in \mathbb{Z}} \mapsto (w_{j+1})_{j \in \mathbb{Z}}$.
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- ▶ Ω always contains at least one element fixed by some power of σ . But observe that not each bi-infinite word fixed by some power of σ is contained in Ω .
- ▶ A substitution dynamical system is minimal and uniquely ergodic.

Desubstitution

Theorem (cf. Mossé (1992))

Let (Ω, S) be the substitution dynamical system induced by σ and $w \in \Omega$. Then there exists a uniquely determined $w' = (w'_j)_{j \in \mathbb{Z}} \in \Omega$ and a uniquely determined integer $k \in \{0, \dots, |\sigma(w'_0)| - 1\}$ such that

$$w = S^k \sigma(w').$$

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By successively applying the theorem on $w = w^{(0)} \in \Omega$ we obtain sequences $(w^{(n)})_{n \geq 1} \in \Omega^{\mathbb{N}}$ and $(k_n)_{n \geq 1} \in \mathbb{Z}^{\mathbb{N}}$ with $0 \leq k_n < |\sigma(w_0^{(n)})|$ for all $n \geq 1$ such that $w^{(n-1)} = S^{k_n} \sigma(w^{(n)})$ (where $w_0^{(n)}$ denotes the letter indexed by 0).

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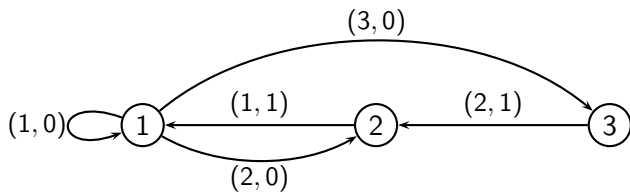
Define

$$\Gamma(w) := (w_0^{(n)}, k_n) \in (\mathcal{A} \times \mathbb{Z})^{\mathbb{N}}.$$

The image $\Gamma(w)$ corresponds to an infinite walk on a finite graph.

Desubstitution

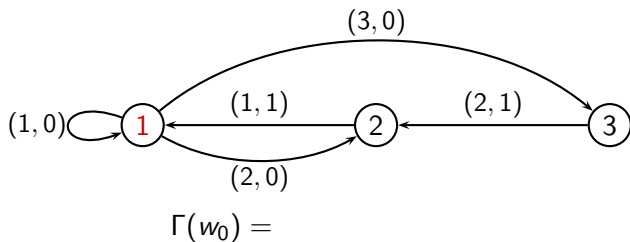
index	...	$\bar{7}$	$\bar{6}$	$\bar{5}$	$\bar{4}$	$\bar{3}$	$\bar{2}$	$\bar{1}$	0	$\dagger 1$	$\dagger 2$	$\dagger 3$	$\dagger 4$	$\dagger 5$	$\dagger 6$	$\dagger 7$...
$w^{(0)} =$...	1	3	1	2	1	1	2	1	3	1	2	1	2	1	3	...



$\Gamma(w_0) =$

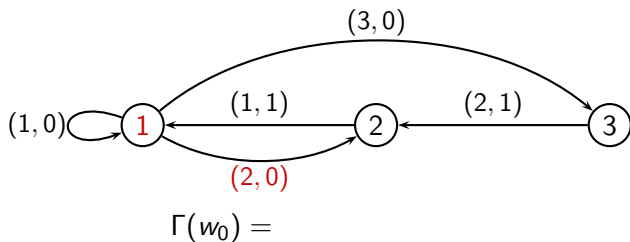
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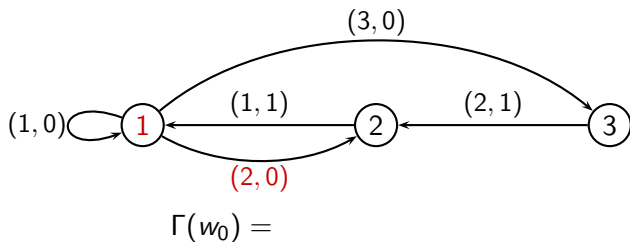
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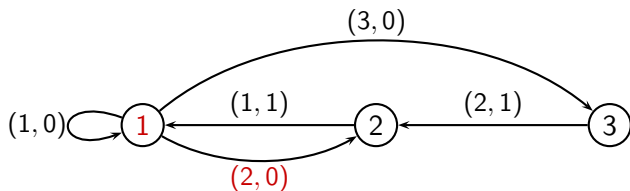
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$w^{(1)} =$								3	1	2	1	1						

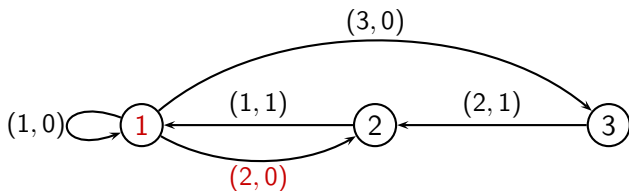


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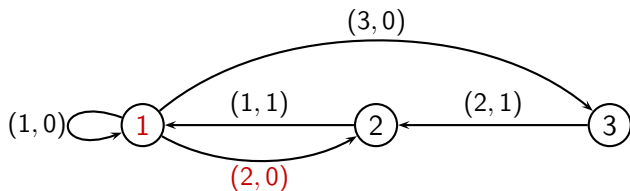
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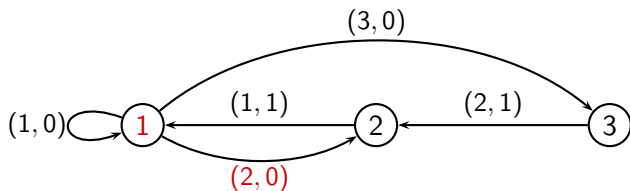
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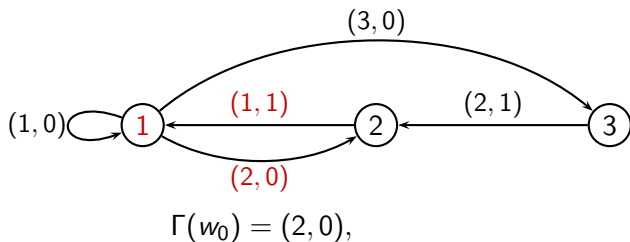
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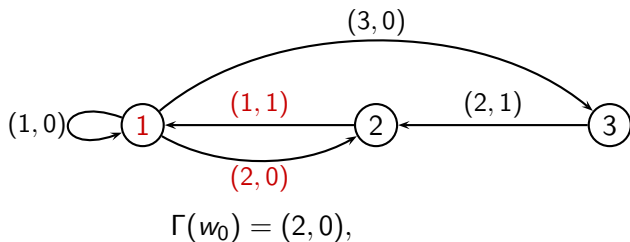
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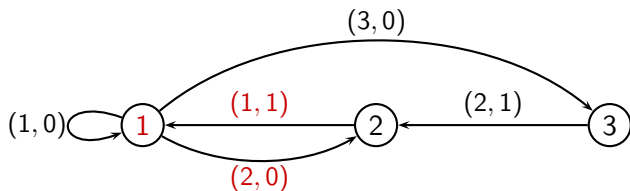
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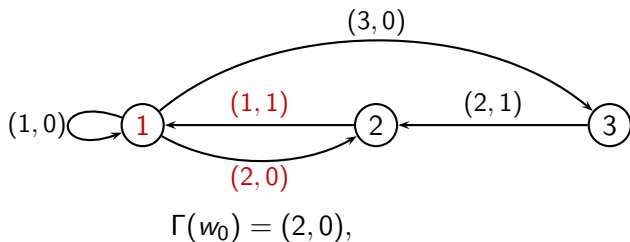
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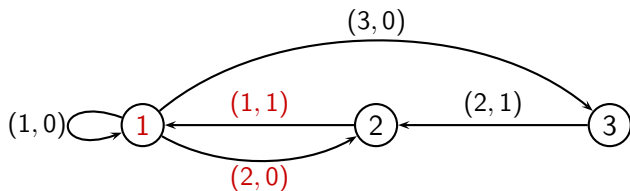
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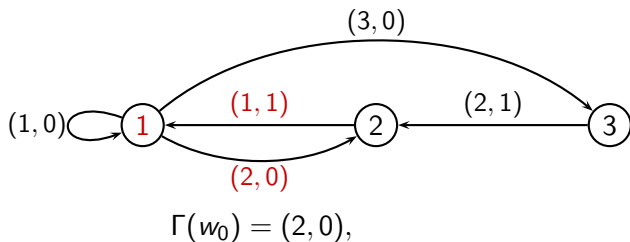
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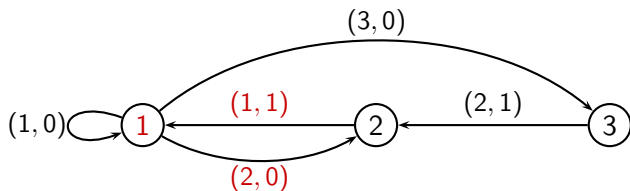
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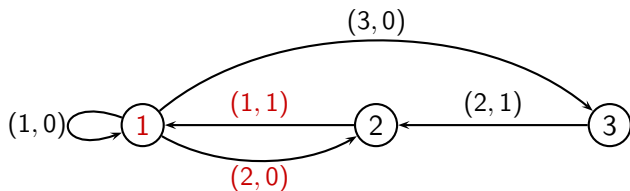
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\vdots									\vdots								



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Shift of finite type

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Shift of finite type

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Theorem (Canterini-Siegel (2001), Holton-Zamboni (2001))

The map Γ is continuous, surjective and, up to a countable set of possible exceptions, injective. The shift S is conjugate to an adic transformation on the edges (Vershik-map).

Scenario

For $w = (w_j)_{j \in \mathbb{Z}} \in \Omega$ let us

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$w^{(0)} =$...	1	3	1	2	1	1		2 . 1		3	1	2	1	2	1	3	...

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$w^{(0)} =$...	1	3	1	2	1	1		2 . 1		3	1	2	1	2	1	3	...
									$k = 0$									
$w^{(1)} =$							3		1 . 2		1							

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$w^{(0)} =$...	1	3	1	2	1	1	2	.	1	3	1	2	1	2	1	3	...
$w^{(1)} =$...	2	1	1	2	1	3	1	.	2	1	1	2	1	3	1	2	...

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$w^{(0)} =$...	1	3	1	2	1	1		2	.	1		3	1	2	1	2	1	3	...	
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$w^{(1)} =$...	2	1	1	2	1	3		1	.	2		1	1	2	1	3	1	2	...	

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$w^{(0)} =$...	1	3	1	2	1	1		2 . 1		3	1	2	1	2	1	3	...
									$k = 0$									
$w^{(1)} =$...	2	1	1	2	1	3		2 . 1		1	1	2	1	3	1	2	...
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$w^{(0)} =$...	1	3	1	2	1	1	2 . 1	3	1	2	1	2	1	3	...
								$k = 0$								
$w^{(1)} =$...	2	1	1	2	1	3	1 . 2	1	1	2	1	3	1	2	...
								$1 \mid 2 . 1 \mid 3$								
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								$k = 0$								
$w^{(1)} =$...	2	1	1	2	1	3	1 . 2 . 1	1	2	1	3	1	2	...	
								...	2 . 1	3	1	2	1	...	($k = 1$)	
								1 . 3							($k = -1$)	

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		...	1	3	1	2	1	1 . 3	1	2	1	1	...	$(k = -1)$		

Question:

What can we do to obtain uniqueness?

Coding prescriptions

Definition

Let c be a function that assigns to each $a \in \mathcal{A}$ a set $c(a) \subset \mathbb{Z}$ that is a complete set of representatives modulo $|\sigma(a)|$ whose absolute values are smaller than $|\sigma(a)|$. (Thus, $|c(a)| = |\sigma(a)|$ and for each $k \in c(a)$ we have $-|\sigma(a)| < k < |\sigma(a)|$; $k' \in c(a)$ with $k' \neq k$ implies $|\sigma(a)| \nmid (k - k')$.) Extend c to \mathcal{A}^2 by defining

$$c(ab) = \{k \in c(a) : k \leq 0\} \cup \{k \in c(b) : k \geq 0\}.$$

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Theorem (S. (2016))

Let (Ω, S) be the substitution dynamical system induced by σ , c a coding prescription, and $w \in \Omega$. Then there exists a uniquely determined $w' = (w'_j)_{j \in \mathbb{Z}} \in \Omega$ and a uniquely determined integer $k \in c(w'_{-1}w'_0)$ such that

$$w = S^k \sigma(w').$$

Coding graphs

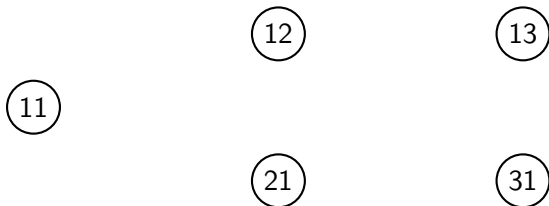
A coding prescription induces in a natural way a graph G_c with set of vertices $\mathfrak{L}_2 := \mathfrak{L} \cap \mathcal{A}^2$ and an edge from $a'b'$ to ab if there is a $k \in c(ab)$ such that $a'b'$ appears at the $(|\sigma(a)| + k)$ th position in $\sigma(ab)$. We label it by (ab, k) .

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$$\sigma(11) = \begin{array}{cccc} & -1 & 0 & 1 \\ 1 & 2 & 1 & 2 \end{array}$$

11

12

13

21

31

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$$\begin{aligned} \sigma(11) &= \overset{-1}{1} \overset{0}{2} \overset{1}{1} \overset{1}{2} \\ c(11) &= \{0, 1\} \end{aligned}$$

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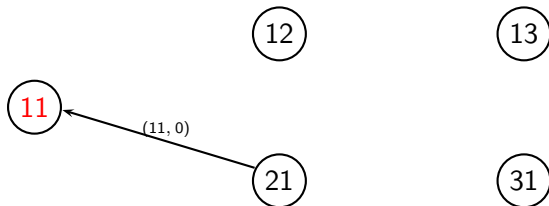
21

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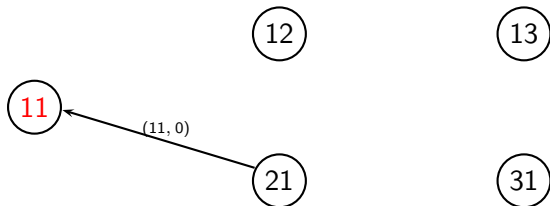
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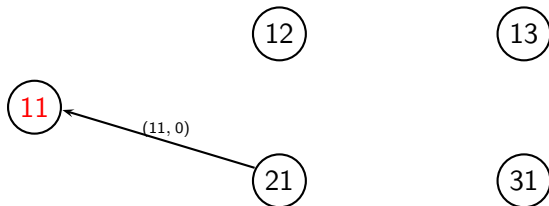
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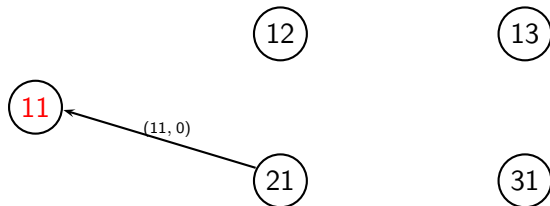


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The set of vertices is $\mathfrak{L}_2 = \{11, 12, 13, 21, 31\}$.

$$\begin{aligned} \sigma(11) &= 1 \quad 2 \quad \overset{-1}{1} \quad \overset{0}{2} \quad \overset{1}{1} \quad \overset{1}{2} \\ c(11) &= \{0, 1\} \end{aligned}$$

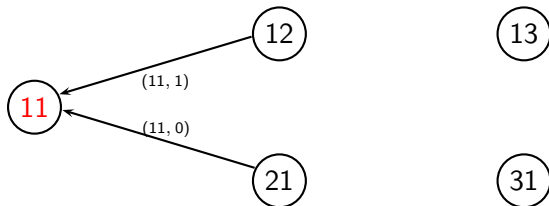


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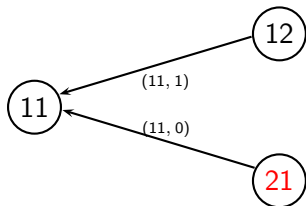
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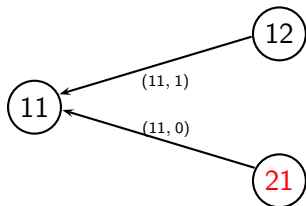
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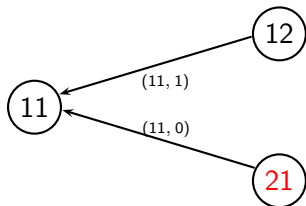


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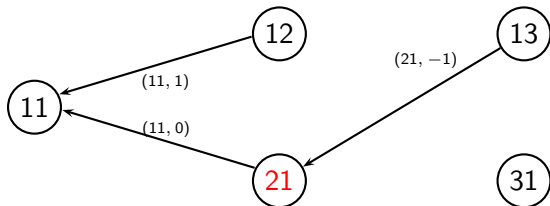
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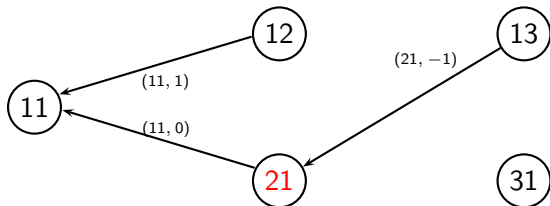
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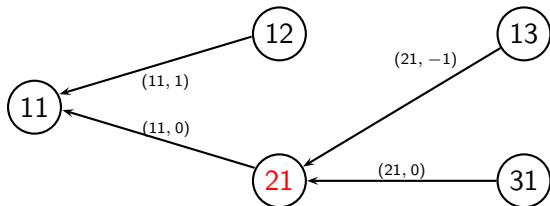
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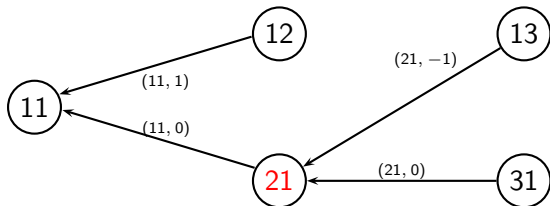
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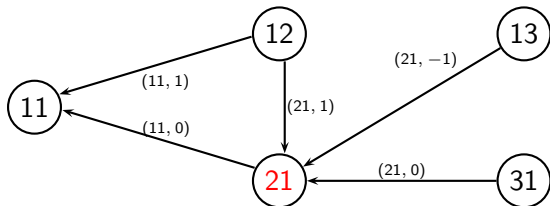
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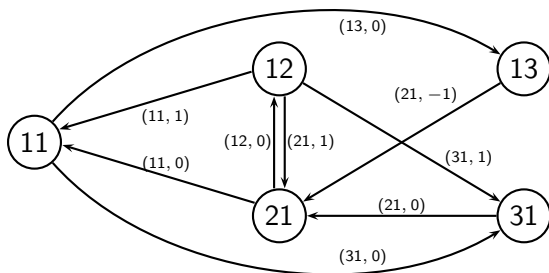
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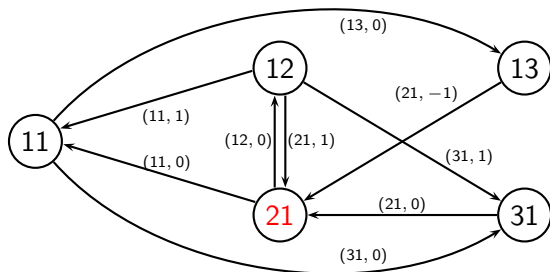
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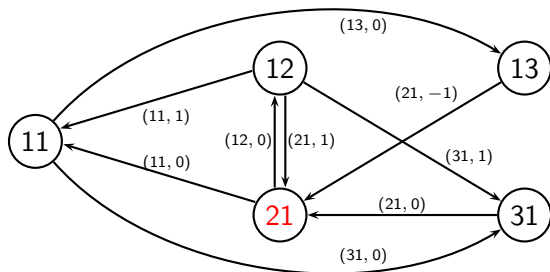
Back to our scenario

index	...	$\bar{7}$	$\bar{6}$	$\bar{5}$	$\bar{4}$	$\bar{3}$	$\bar{2}$	$\bar{1}$	\cdot	0	\dagger	\dagger	\dagger	\dagger	\dagger	\dagger	...	
$w^{(0)} =$...	1	3	1	2	1	1	2	.	1	3	1	2	1	2	1	3	...



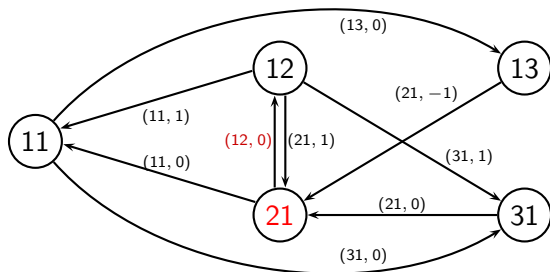
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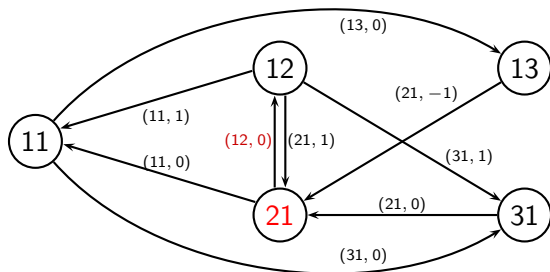
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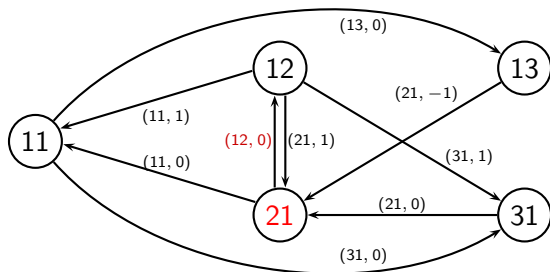
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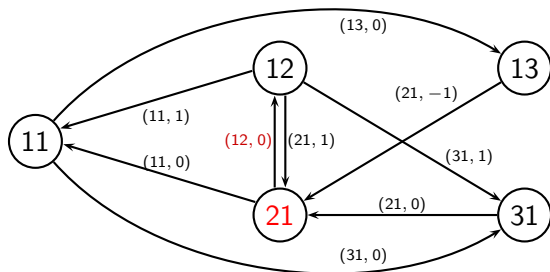
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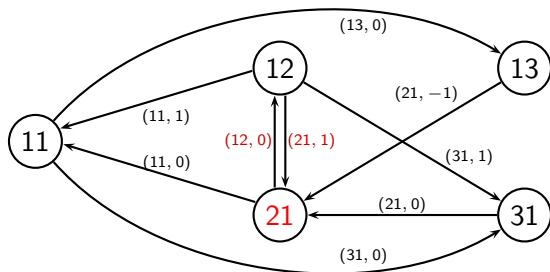
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Shifts of finite type

Let $(w^{(n)})_{n \geq 1} \in \Omega^{\mathbb{N}}$ and $(k_n)_{n \geq 1} \in \mathbb{Z}^{\mathbb{N}}$ with $k_n \in c(w_{-1}^{(n)} w_0^{(n)})$ for all $n \geq 1$ (where $w_{-1}^{(n)} w_0^{(n)}$ denotes central pair of $w^{(n)}$) such that $w^{(n-1)} = S^{k_n} \sigma(w^{(n)})$. Define

$$\Gamma_c(w) := (w_{-1}^{(n)} w_0^{(n)}, k_n) \in (\mathfrak{L}^2 \times \mathbb{Z})^{\mathbb{N}}.$$

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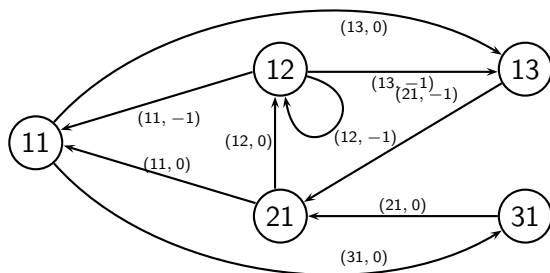
Theorem (S. (2016))

The map Γ_c is continuous, surjective and, up to a countable set of possible exceptions, injective.

Comparison

For each substitution σ there are exactly $\prod_{a \in \mathcal{A}} 2^{|\sigma(a)|-1}$ different coding prescriptions.

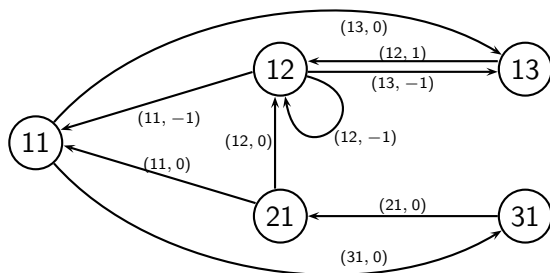
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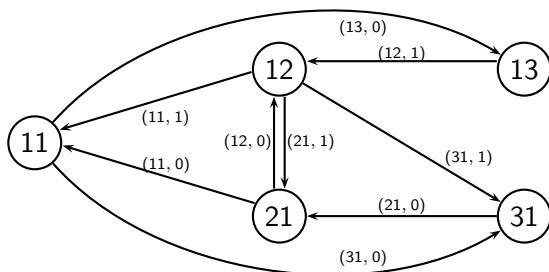
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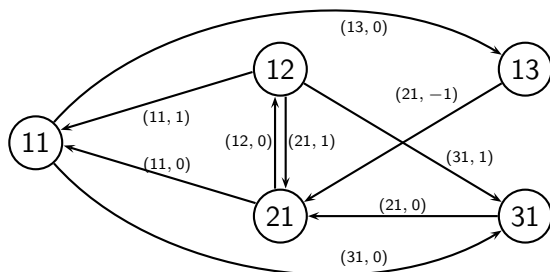
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- ▶ We always can find a coding prescription such that Γ_c is completely injective (possibly by considering higher powers of σ).
- ▶ Choose c such that $c(a)$ contains only non-negative integers for all $a \in \mathcal{A}$ and $w \in \Omega$. If $\Gamma_c(w) = (a_n b_n, k_n)_{n \geq 1}$ then $\Gamma(w) = (b_n, k_n)_{n \geq 1}$. Thus, the coding mentioned at the beginning corresponds to one special choice of c .

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We call a coding prescription *continuous* if $c(a)$ consists of consecutive integers for each $a \in \mathcal{A}$.

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Theorem (S. (2016))

If c is continuous then Γ_c conjugates the shift map S with an adic transformation on the edges.

Application: Numeration

Currently I work on an application of the theory in numeration. For a given primitive substitution we can associate with each coding prescription a number system. The coding prescription that assigns to each letter a set of non-negative integers yields the well-known Dumont-Thomas numeration.

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More details: 2017 in Rome

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Thank you for your attention